

# SODMRP: Stable On-Demand Multicast Routing Protocol

**Abstract.** A mobile ad hoc network (MANET) is a dynamically reconfigurable wireless network that does not have a fixed infrastructure. The group-oriented services are one of the important application classes that are adopted by mobile ad hoc networks (MANETs). To support such services, multicast routing is used. But, stable multicast routing is difficult in MANETs due to limited resources and mobility of nodes. This paper proposes an energy-aware approach to develop a stable multicast routing protocol. By offering some modifications over ODMRP protocol, we involve nodes' energy information in the proposed protocol, to increase routes stability. Simulation results show that the proposed protocol improves route stability and hence leads to more packet delivery rate (PDR), more throughput and longer route life time, in compare with ODMRP.

**Streszczenie.** Mobilne bezprzewodowe sieci ad-hoc MANET nie mają stałej infrastruktury. W artykule zaproponowano stabilny protokół przesyłania grupowego informacji umożliwiający zwiększenie szybkości komunikacji pakietowej. (**SODMRP – stabilny protokół grupowego przesyłania informacji w bezprzewodowych sieciach mobilnych**)

**Keywords:** MANET, Multicast, ODMRP, Stability, Energy Consumption.

**Słowa kluczowe:** MANET, sieci bezprzewodowe, sieci mobilne, routing protocol – protokół trasowania.

## Introduction

In the application level, MANET users communicate with together as teams. Thus applications need for group communication (multicasting) to data forwarding and real time traffics. With the rapid growth of group communication services, the multicast routing in MANET has attracted a lot of attention recently [1][2][3][4][5]. Existing protocols are either tree or mesh-based. Examples of tree-based schemes include [6][7][8]: Ad hoc multicast routing protocol (AMRoute), ad hoc multicast routing utilizing increasing ID-numbers protocol (AMRIS), and multicast ad hoc on-demand distance vector routing protocol (MAODV). On the other hand, some examples of mesh-based schemes are (a) on demand multicast routing protocol (ODMRP [9]), (b) forwarding group multicast protocol (FGMP[10]), (c) core assisted mesh protocol (CAMP[11]), (d) neighbor supporting ad hoc multicast routing protocol (NSMP[12]), (e) location-based multicast protocol[13], and (f) dynamic core-based multicast protocol (DCMP[14]). ODMRP is one of the well established protocols, whose major advantage is its high packet delivery ratio and throughput even under highly mobile network conditions. The disadvantage of ODMRP is that the control overhead grows higher and higher with network size [15][16][17][18]. Another important issue that has attracted remarkable attention is route stability. Because of this importance, [19] constructs a new metric of node stability and selects a stable path by using entropy metric to reduce the number of route reconstruction.

This paper provides an extension to ODMRP protocol that makes it more stable by using an energy-aware approach. While ODMRP uses only movement information to provide route stability, the proposed approach considers also nodes' residue energy to increase the route stability.

## ODMRP Protocol

ODMRP [9] is a mesh-based rather than a conventional tree-based scheme and uses a forwarding group concept. ODMRP uses two types of packet to establish and update routes: Join-Query and Join-Reply. We can summarize route discovery phase of ODMRP as follows:

1. Multicast source originates a Join-Query packet. When location and movement information is utilized, it sets the MIN LET (Link Expiration Time) field of the packet to the MAX LET VALUE.

2. When a node receives a Join-Query packet, checks if it is a duplicate. If duplicate, then discard the packet.

3. If it is not a duplicate, insert an entry into message cache with the information of the received packet (i.e., sequence number and source IP address) and insert/update the entry for routing table (i.e., backward learning).

4. Use movement information to predict the duration of time the link between the node and the upstream node will remain connected. Assume node  $i$  is the upstream node and node  $j$  is the current node. Let  $(x_i, y_i)$  be the coordinate of node  $i$  and  $(x_j, y_j)$  be that of node  $j$ . Also let  $v_i$  and  $v_j$  be the speeds, and  $\theta_i$  and  $\theta_j$  be the moving directions of nodes  $i$  and  $j$ , respectively. The information of node  $i$  can be obtained from the Join-Query and the current node's location and mobility information can be provided by the GPS. The duration of time that the link between two nodes will stay connected, i.e. LET, is given by equation (1).

$$(1) \quad LET = \frac{-(ab + cd) + \sqrt{(a^2 + c^2)r^2 - (ad - cd)^2}}{(a^2 + c^2)} \quad (1)$$

where:

$$\begin{aligned} c &= v_i \sin \theta_i - v_j \sin \theta_j, \quad d = y_i - y_j, \\ b &= x_i - x_j, \quad a = v_i \cos \theta_i - v_j \cos \theta_j \end{aligned}$$

The minimum between this LET value and the indicated value in MIN LET field of the Join-Query is included in the packet. Minimum LET in any route is known as RET (Route Expiration Time).

5. The receiver then chooses the most stable route (i.e., the route with the largest RET) and transmits a Join-Reply packet.

More details about ODMRP can be found in [9].

## Proposed Protocol: SODMRP

To more carefully predict of rout expiration time and hence improve route stability, we include residue energy of nodes in our protocol. As the first step, let's define a parameter that plays important role in our work, namely, Residue energy (RES). For any node, RES is the amount of remaining energy of the node. To update RES, any node uses  $E_{tx}$  and  $E_{rx}$  that are amounts of energy required to transmit and to receive of a packet. Thus, node energy consumption (NEC) of the node is computed as follows for

transmission of  $n$  packets (original and retransmitted packets) by energy module:

$$(2) \quad NEC = n * (E_{tx} + E_{rx})$$

Now, we can update  $RES$  of the node by using equation (3).

$$(3) \quad RES_{new} = RES_{old} - NEC$$

As the first requirement of this protocol, we add a new field called MIN\_RES to Join-Query packet header, to carry minimum RES of the routes. Now, we can describe main operations of SODMRP as follows:

1. At the beginning, multicast source originates a Join Query packet and sets the MIN LET and MIN RES fields to the MAX LET VALUE and MAX RES VALUE, respectively.

2. When a node receives a Join-Query packet, checks if it is a duplicate. If yes, then discards it.

3. If it is not a duplicate, inserts an entry into message cache with the information of the received packet.

4. Updates MIN LET and MIN RES fields of received Join-Query packet:

4. a. Predicts the duration of time the link between the node and the upstream node will remain connected (LET), same as ODMRP, by using equation (1). The minimum between this LET value and the indicated value in MIN LET field of the Join-Query is included in the packet.

4. b. Computes RES value and includes the minimum between this RES value and the value in MIN RES field of the Join-Query in the packet.

5. To select a route, the receiver chooses the most stable route which is selected based on the weighted equation of (4), in which,  $MaxRET_k$  and  $MaxRES_k$  are the largest MaxRET and MaxRES received at source k. Each route i has its own  $RET_i$  and  $RES_i$ , hence receiver can compute  $W_i$  by using equation (4). Finally receiver selects the most stable route (i.e., the route with the largest  $W_i$ ) and then transmits a Join Reply packet.

$$(4) \quad W_i = C_1 * \left[ \frac{RET_i}{MaxRET_k} \right] + C_2 * \left[ \frac{RES_i}{MaxRES_k} \right]$$

In equation (4) constants  $C_1$  and  $C_2$  represent the weights of RET and RES.  $C_1$  and  $C_2$  are positive constants, and we take  $C_1+C_2=1$ . The values of  $C_1$  and  $C_2$  should be chosen according to the system requirements. Suitable selection of  $C_1$  and  $C_2$  provides a better QoS for multicast applications. For example in military environment energy may be more important than mobility, therefore  $C_2$  should be selected larger than  $C_1$  e.g.  $C_1=0.3$  and  $C_2=0.7$ . In this paper we use  $C_1=C_2=0.5$  to give same importance to energy and mobility information.

Other operations of SODMRP are same as ODMRP protocol.

### Evaluation of ODMRP Protocol

In this section SODMRP is evaluated in terms of design necessity and performance.

#### a. Evaluation of ODMRP Protocol

Note that there are many routing protocols for MANETs and ODMRP is one of them that offers high throughput even in highly mobile network conditions. Naturally, each protocol has some advantages and also suffers from disadvantages. In ODMRP protocol, instability of routes is an important weakness that results in links breakage and leads in turn to low performance and high control overhead

of network. This work aims to design an efficient routing algorithm for MANETs by considering this issue. It can be found in Tab.1 how this work differs from other protocols derived from ODMRP.

Table 1. SODMRP and other protocols derived from ODMRP [16,20-26]

| Protocols | Main Idea  |
|-----------|--|
| ODMRP-LR  | Link failure detection and recovery  |
| RODMRP    | Offers more reliable forwarding paths in face of node failure and network failures                           |
| SC-ODMRP  | Tries to increase scalability of ODMRP   |
| CODMRP    | Attempts to improve PDR efficiency by minimizing the data redundancy   |
| PSO-ODMRP | Uses nature principle exists in the form of fitness function (PSO algorithm)                                 |
| IODMRP    | Improves ODMRP by decreasing control overhead in case of increased multicast group                           |
| EODMRP    | Uses motion adaptive refresh technique for route local recovery to achieve minimum overhead                  |
| RDTODMRP  | Cuts down the unnecessary redundant routes and their data transmissions for obtaining better result          |
| E-ODMRP   | Use local recovery technique for decrease link breakage  |
| SODMRP    | Use residue energy of nodes for obtaining route expire time for select stable route by route weight function |

#### b. Performance Evaluation

To investigate the impact of our enhancements, we implement SODMRP by making some modifications on ODMRP module of GloMoSim [27] software package. A total of 100 nodes are simulated for duration of 600 second in an area of 1000m×1000m. The mobility model is the random way point to model the mobility of the nodes in the network with the pause time of 10 ms. The MAC layer protocol used is IEEE 802.11. The transmission range for each node is 250 m and the channel capacity is 2 Mbps. The initial energy of the node is 1500 joules. We simulate this network once under ODMRP protocol and then under SODMRP protocol to compare their performance. The metrics used for comparison are packet delivery, throughput and life time of the selected routes. The performance of the protocols and evaluation results are illustrated versus mobility speed and number of nodes. We consider two set of experiments for both SODMRP and ODMRP. In the first set of experiments, the mobility speed is set constant to 20 m/h and number of mobile nodes is varied from 10 to 100. In the second experiment the number of nodes is set to 100 and mobility speed is varied from 10 to 70 m/h.

Figs. 1 and 2 shows the simulation results. Fig.1 shows that as the number of nodes increase (from 10 to 100), there is an increase in route life time, PDR and throughput, for both ODMRP and SODMRP. This is reasonable, since there is a possibility of decreased hop length between the nodes with increase in number of nodes, which enhances the link stability and reduces link breakages and packet drops. On the other hand, according to this figure SODMRP behaves better than ODMRP in terms of throughput and packet delivery rate for various numbers of nodes. This is due to the higher stability of SODMRP routes that reduces packet loss in case of link or node failures. Fig. 2 on the other hand, shows how SODMRP outperforms ODMRP for a wide range of mobility speed. As the previous experiment, because of the higher stability, SODMRP experiences less route breakage and hence achieves better performance in compare with ODMRP.

As another step in performance evaluation of SODMRP, Tab. 2 shows its average behavior in compare with other ODMRP-based routing protocols. This table covers those protocols that been evaluated in terms of throughput, packet delivery rate and route life time in corresponding

papers [20-24]. It shows that improvement percentage made by SODMRP over ODMRP is remarkably more than other ones.

Table 2. Average improvement percentage over ODMRP for various protocols

| Parameters | RDTODMR | EODMRP | SC-ODMRP | SODMRP |
|------------|---------|--------|----------|--------|
| Throughput | 7%      | 8%     | 11%      | 14%    |
| PDR        | 5%      | 7%     | 9%       | 13.5%  |
| Lifetime   | 2.5%    | 4.5%   | 9.5%     | 15%    |

## Conclusion

This paper proposed a stable on-demand multicast routing scheme in MANET. The scheme finds stable multicast routes to receivers by considering node's residue energy. It uses a route weighted function in ODMRP route discovery process, to consider both movement and energy information in its operation. Simulation results illustrate that this approach leads to better performance in terms of packet delivery rate, throughput and route life time, in compare with ODMRP and other ODMRP-based protocols.

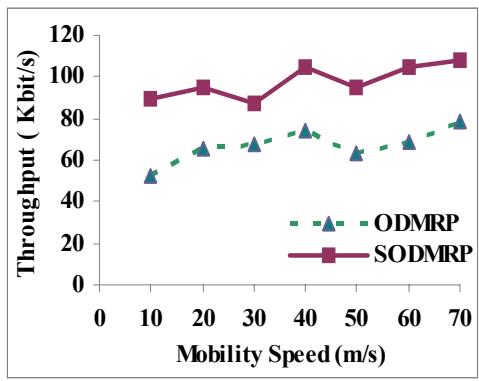
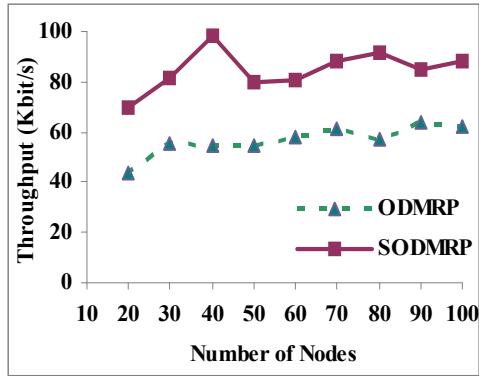
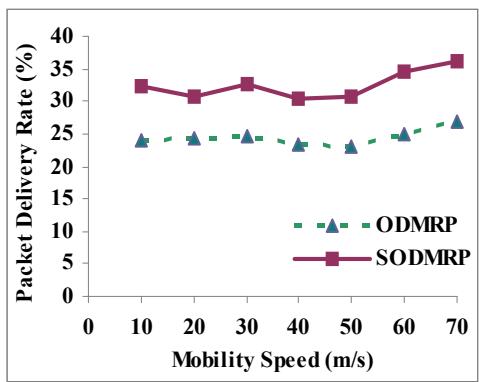
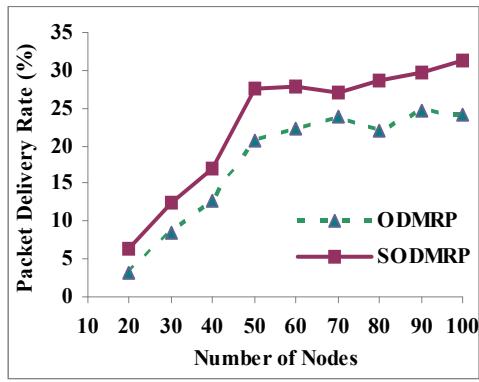
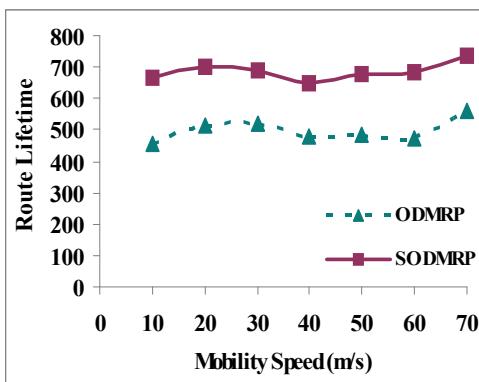
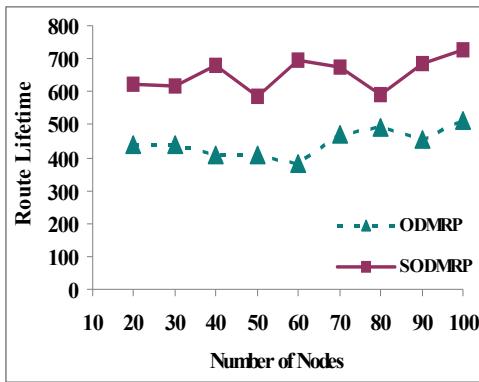


Fig.1. Experiment 1: Mobility speed=20 m/h.

Fig.2. Experiment 2: Number of nodes=100.

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